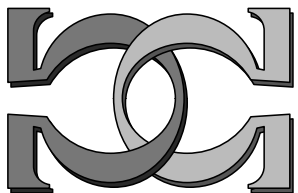
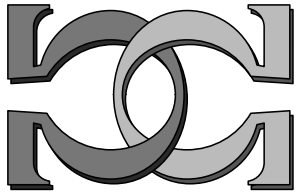
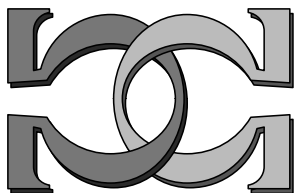


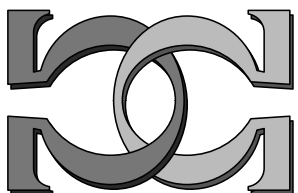
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in Physics**



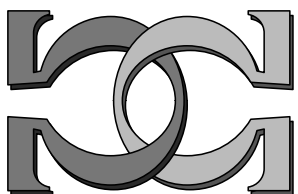
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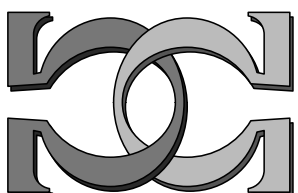
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Article

A Non-Probabilistic Model of Relativised Predictability in Physics

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Abstract: Little effort has been devoted to studying generalised notions or models of (un)predictability, yet is an important concept throughout physics and plays a central role in quantum information theory, where key results rely on the supposed inherent unpredictability of measurement outcomes. In this paper we continue the programme started in [1] developing a general, non-probabilistic model of (un)predictability in physics. We present a more refined model that is capable of studying different degrees of “relativised” unpredictability. This model is based on the ability for an agent, acting via uniform, effective means, to predict correctly and reproducibly the outcome of an experiment using finite information extracted from the environment. We use this model to study further the degree of unpredictability certified by different quantum phenomena, showing that quantum complementarity guarantees a form of relativised unpredictability that is weaker than that guaranteed by Kochen-Specker-type value indefiniteness. We exemplify further the difference between certification by complementarity and value indefiniteness by showing that, unlike value indefiniteness, complementarity is compatible with the production of computable sequences of bits.

Keywords: prediction; unpredictability; randomness; complementarity

17 **1. Introduction**

18 Many physical processes and phenomena are intuitively thought of as unpredictable: the roll of a
19 die, the evolution of weather systems, and the outcomes of quantum measurements, to mention a few.
20 While *ad hoc* definitions of unpredictability may exist within certain domains, little work has been done
21 towards developing a more general understanding of the concept. Although domain specific notions
22 of unpredictability may help describe and categorise phenomena within the domain, the concept of
23 unpredictability has a much more central and important role in quantum information theory.

24 Many of the advantages promised by quantum information theory and cryptography rely critically
25 on the belief that the outcomes of quantum measurements are intrinsically unpredictable [2,3]. This
26 belief underlies the use of quantum random number generators to produce “quantum random” sequences
27 that are truly unpredictable (unlike pseudo-randomness) [4] and the generation of cryptographic keys
28 unpredictable to any adversary [3]. Such claims of quantum unpredictability are generally based
29 on deeper theoretical results—such as the Kochen-Specker [5] and Bell [6] theorems, or quantum
30 complementarity—but nonetheless remain informal intuition.

31 The quantum cryptography community has used a probability theoretic approach to try and make
32 use of, and quantify the degree of unpredictability in quantum information theoretical situations, in
33 particular by following the cryptographic paradigm of adversaries with limited side-information [7].
34 This approach, while suitable in such cryptographic situations precisely because of its epistemic
35 nature [8], relies on the probabilistic formalism of quantum mechanics and the subsequently assumed
36 unpredictability. In order to fully understand and study the degree of quantum unpredictability and
37 randomness, it is instead crucial to have more general models of unpredictability to apply.

38 Historically, little work has been devoted to such generalised notions of unpredictability. In [1] we
39 discussed in some detail the most notable approaches, in particular those of Popper [9], Wolpert [10],
40 and Eagle [11]. In response to these approaches, we outlined a new model based around the ability for a
41 predicting agent, acting via uniform, effective means, to predict correctly and reproducibly the outcome
42 of an experiment using some finite information the agent extracts from the “environment” as input.

43 This model allowed us to consider a specific, ontic, form of unpredictability which was particularly
44 suited for analysing the type of unpredictability quantum mechanics claims to provide. However, this
45 strong form of unpredictability is too strong in many cases and failed to capture the possible different
46 degrees of unpredictability: what is predictable for one agent may not be for another with different
47 capabilities.

48 In this paper we refine and improve this model of (un)predictability, providing a more nuanced,
49 relativised notion of unpredictability that can take into account the epistemic limits of an observer,
50 something crucial, for example, in chaotic systems [12]. This also provides the ability to look at the
51 degree of unpredictability guaranteed by different possible origins of quantum unpredictability. We
52 examine one such case—that of quantum complementarity—in detail, and show that it provides a weaker
53 form of unpredictability than that arising from Kochen-Specker-type value indefiniteness as discussed
54 in [1]

55 **2. Relativised model of predictability**

56 The model of (un)predictability that we proposed in [1] is based around the ability of an agent to, in
57 principle, predict the outcome of a physical experiment. By using computability theory—motivated
58 by the Church-Turing thesis—to provide a universal framework in which prediction can occur,
59 this information-theoretical approach allows different physical systems and theories to be uniformly
60 analysed.

61 Here we refine and extend this model to be able to relativise it with respect to the means/resources of
62 the predicting agent. This gives our model an epistemic element, where our previous and more objective
63 model can be obtained as the limit case. In this framework we can consider the predictive capabilities
64 of an agent with limited capacities imposed by practical limitations, or under the constraints of physical
65 hypotheses restricting such abilities.

66 Before we proceed to present our model in detail, we will briefly outline the key elements comprising
67 it.

- 68 1. The specification of an experiment E for which the outcome must be predicted.
- 69 2. A predicting agent or “predictor”, which must predict the outcome of the experiment. We model
70 this as an effectively computable function, a choice which we will justify further.
- 71 3. An extractor ξ is a physical device the agent uses to (uniformly) extract information pertinent to
72 prediction that may be outside the scope of the experimental specification E . This could be, for
73 example, the time, measurement of some parameter, iteration of the experiment, etc.
- 74 4. A prediction made by the agent with access to a set Ξ of extractors. The set of extractors Ξ provides
75 the relativisation of the model.

76 This model is explicitly a non-probabilistic one, a fact that may seem overly restrictive given that
77 highly probable events seem predictable. However, the uncertainty present in “high probabilities”
78 represents an important latent unpredictability in such processes, and certainty is needed if predictions
79 are to be related to definite properties of physical systems [13], as in quantum scenarios, for example.

80 It should be noted that our model does not assess the ability to make statistical predictions about
81 physical processes (as one might about the throw of a dice, for example)—as probabilistic models
82 might—but rather the ability to predict precise measurement outcomes.

83 We will next elaborate on the individual aspects of the model.

84 2.1. Predictability model

85 **Experimental specification.** An experiment is a finite specification for which the outcome is to
86 be predicted. We restrict ourselves to the case where the result of the experiment, i.e. the value to be
87 predicted, is a single bit: 0 or 1. However, this can readily be generalised for any finite outcome. On the
88 other hand it does not make sense to predict an outcome requiring an infinite description, such as a real
89 number, since this can never be measured exactly. In such a case the outcome would be an approximation
90 of the real—a rational number, and thus finitely specifiable.

91 The experimental specification, being finite, cannot normally specify exactly the required setup
92 of the experiment, as a precise description of experimental conditions generally involves real-valued

93 parameters. Rather, it is expressed with finite precision by the experimenter within their limited
94 capacities—making use, for example, of the pertinent symmetries to describe the experiment. A
95 particular trial of E is associated with the parameter λ which fully describes the “state of the universe”
96 in which the trial is run. As an example, one could consider E to specify the flipping of a certain coin,
97 or it could go further and specify, up to a certain accuracy, the initial dynamical conditions of the coin
98 flip. In both cases, λ contains further details—such as the exact initial conditions—which could be used
99 by an agent in trying to predict the result of E .

100 The parameter λ will generally¹ be “an infinite quantity”—for example, an infinite sequence or a
101 real number—structured in an unknown manner. Forcing a specific encoding upon λ , such as a real
102 number, may impose an inadequate structure (e.g. metric, topological) which is not needed for what
103 follows. While λ is generally not in its entirety an obtainable quantity, it contains any information that
104 may be pertinent to prediction—such as the time at which the experiment takes place, the precise initial
105 state, any hidden parameters, etc.—and any predictor can have practical access to a finite amount of this
106 information. We can view λ as a resource from which one can extract finite information in order to try
107 and predict the outcome of the experiment E .

108 **Predicting agent.** The predicting agent (or “predictor”) is, as one might expect, the agent trying to
109 predict the outcome of a particular experiment, using potentially some data obtained from the system
110 (i.e. from λ) to help in the process. Since such an agent should be able to produce a prediction in a finite
111 amount of time via some uniform procedure, we need the prediction to be *effective*.

112 Formally, we describe a predicting agent as a computable function P_E (i.e. an algorithm) which halts
113 on every input and outputs either 0,1, or “prediction withheld”. Thus, the agent may refrain from making
114 a prediction in some cases if it is not certain of the outcome. P_E will generally be dependent on E ,
115 but its definition as an abstract algorithm means *it must be able to operate without interacting with the*
116 *subsystem whose behaviour it predicts*. This is necessary to avoid the possibility that the predictor affects
117 the very outcome it is trying to predict.

118 We note finally that the choice of computability as the level of effectivity required can be strengthened
119 or weakened, as long as some effectivity is kept. Our choice of computability is motivated by the
120 Church-Turing thesis, a rather robust assumption [14].

121 **Extractor.** An extractor is a physically realisable device which a predicting agent can use to extract
122 (finite) useful data that may not be a part of the description of E from λ to use for prediction—i.e. as
123 input to P_E . In many cases this can be viewed as a measurement instrument, whether it be a ruler, a
124 clock, or a more complicated device.

125 Formally, an extractor produces a finite string of bits $\xi(\lambda)$ which can be physically realised without
126 altering the system, i.e. passively. In order to be used by P_E for prediction, $\xi(\lambda)$ should be finite and
127 effectively codable, e.g. as a binary string or a rational number.

¹ If one insists on a discrete or computational universe—whether it be as a “toy” universe, in reality or in virtual reality—then λ could be conceived as a finite quantity. This is, however, the exception, and in the orthodox view of real physical experiments λ would be infinite, even if the prediction itself is discrete or finite, so we will adopt this view here.

128 **Prediction.** We define now the notion of a correct prediction for a predicting agent having access to
 129 a fixed (finite or infinite) set Ξ of extractors.

130 Given a particular extractor ξ , we say the prediction of a run of E with parameter λ is *correct for ξ* if
 131 the output $P_E(\xi(\lambda))$ is the same as the outcome of the experiment. That is, it correctly predicts E when
 132 using information extracted from λ by ξ as input.

133 However, this is not enough to give us a robust definition of predictability, since for any given run
 134 it could be that we predict correctly by chance. To overcome this possibility, we need to consider the
 135 behaviour of repeated runs of predictions.

136 A *repetition procedure for E* is an algorithmic procedure for resetting and repeating the experiment
 137 E . Generally this will be of the form “ E is prepared, performed and reset in a specific fashion”. The
 138 specific procedure is of little importance, but the requirement is needed to ensure the way the experiment
 139 is repeated cannot give a predicting agent power that should be beyond their capabilities or introduce
 140 mathematical loopholes by “encoding” the answer in the repetitions; both the prediction and repetition
 141 should be performed algorithmically.

142 We say the predictor P_E is correct for ξ if for any k and any repetition procedure for E (giving
 143 parameters $\lambda_1, \lambda_2, \dots$ when E is repeated) there exists an $n \geq k$ such that after n repetitions of E
 144 producing the outputs x_1, \dots, x_n , the sequence of predictions $P_E(\xi(\lambda_1)), \dots, P_E(\xi(\lambda_n))$:

- 145 1. contains k correct predictions,
- 146 2. contains no incorrect prediction; e.g. the remaining $n - k$ predictions are withheld.

147 From this notion of correctness we can define predictability both relative to a set of extractors, and in
 148 a more absolute form.

149 Let Ξ be a set of extractors. An experiment E is *predictable for Ξ* if there exists a predictor P_E and
 150 an extractor $\xi \in \Xi$ such that P_E is correct for ξ . Otherwise, it is *unpredictable for Ξ* .

151 This means that P_E has access to an extractor $\xi \in \Xi$ which, when using this extractor to provide
 152 input to P_E , can be made to give arbitrarily many correct predictions by repeating E enough (but finitely
 153 many) times, without ever giving an incorrect prediction.

154 The more objective notion proposed in [1] can be recovered by considering all possible extractors.
 155 Specifically, an experiment is (*simply*) *predictable* if there exists a predictor P_E and an extractor ξ such
 156 that P_E is correct for ξ . Otherwise, it is (*simply*) *unpredictable*.

157 The outcome x of an *single trial* of the experiment E performed with parameter λ is *predictable (for*
 158 $\Xi)$ if E is predictable (for Ξ) and $P_E(\xi(\lambda)) = x$. Otherwise, it is *unpredictable (for Ξ)*. We emphasise
 159 here that the predictability of the result of a single trial is *predictability with certainty*.

160 2.2. Relativisation

161 While the notion of simple predictability provides a very strong notion of unpredictability—one that
 162 seems to correspond to what is often meant in the context of quantum measurements [1]—in some
 163 physical situations, particularly in classical physics, our inability to predict would seem to be linked to
 164 our epistemic lack of information, often due to measurement. Put differently, unpredictability is a result

165 of only having access to a set Ξ of extractors of limited power. Our relativised model of prediction
 166 attempts to capture this, defining predictability relative to a given set of extractors Ξ .

167 2.2.1. Specifying the set of extractors Ξ

168 In defining this notion, we deliberately avoided saying anything about how Ξ should be specified.
 169 Here we outline two possible ways this can be done.

170 The simplest, but most restrictive, way would be to explicitly specify the set of extractors. As an
 171 example, let us consider the experiment of firing a cannonball from a cannon and the task of predicting
 172 where it will land (assume for now that the muzzle velocity is known and independent of firing angle).
 173 Clearly, the position will depend on the angle the cannonball is fired at. Then, if we are limited to
 174 measuring this with a ruler, we can consider, for example, the set of extractors

$$\Xi = \{\xi \mid \xi(\lambda) = (x, y) \text{ where } x \text{ and } y \text{ are the muzzle position to an accuracy of } 1\text{cm}\}$$

175 and then consider predictability with respect to this set Ξ . (For example, by using trigonometry to
 176 calculate the angle of firing, and then where the cannonball will land.)

177 Often it is unrealistic to characterise completely the set of extractors available to an agent in this
 178 way—think about a standard laboratory full of measuring devices that can be used in various ways.
 179 Furthermore, such devices might be able to measure properties indirectly, so we might not be able
 180 to characterise the set Ξ so naively. Nonetheless, this can allow simple consideration and analysis of
 181 predictability in various situations, such as under-sensitivity to initial conditions.

182 A more general approach, although often requiring further assumptions, is to limit the “information
 183 content” of extractors. This avoids the difficulty of having to explicitly specify Ξ . Continuing with the
 184 same example as before, we could require that no extractor $\xi \in \Xi$ can allow us to know the firing angle
 185 better than 1° . This circumvents any problems raised by the possibility of indirect measurement, but of
 186 course requires us to have faith in the assumption that this is indeed the case; it could be possible that we
 187 *can* extract the angle better than this, but we simply don’t know how to do it with our equipment. (This
 188 would not be a first in science!) Nonetheless, this approach captures well the epistemic position of the
 189 predicting agent.

190 Let us formalise this more rigorously. We hypothesise that we cannot do any better than a hypothetical
 191 extractor ξ' extracting the desired physical quantity. Then we characterise Ξ by asserting: for all $\xi \in \Xi$
 192 there is no computable function f such that for every parameter λ , $f(\xi(\lambda))$ is more accurate than ξ' .
 193 Obviously, the evaluation of “more accurate” requires a (computable) metric on the physical quantity
 194 extracted, something not unreasonable physically given that observables tend to be measured as rational
 195 numbers as approximations of reals [15].

196 This general approach would need to be applied on a case by case basis, given assumptions about
 197 the capabilities available to the predicting agent. Assumptions have to be carefully justified and, ideally,
 198 subject themselves to experimental verification.

199 Either of these approaches, and perhaps others, can be used with our relativised model of prediction.
 200 In any such case of relativisation, one would need to argue that the set Ξ unpredictability is proven for is
 201 relevant physically. This is unavoidable for any epistemic model of prediction.

202 2.2.2. A detailed example

203 Let us illustrate the use of relativised unpredictability with a more interesting example of an
 204 experiment which is predictable, but its intuitive unpredictability is well captured by the notion of
 205 relativised unpredictability. In particular, let us consider a simple chaotic dynamical system. Chaos is
 206 often considered to be a form of unpredictability, and is characterised by sensitivity to initial conditions
 207 and the mixing of nearby dynamical trajectories [12]. However, chaos is, formally, an asymptotic
 208 property [16], and we will see that as a result the unpredictability of chaotic systems is not so simple as
 209 might be initially suspected.

210 For simplicity, we will take the example of the dyadic map, i.e. the operation on infinite sequences
 211 defined by $d(x_1x_2x_3\dots) = x_2x_3\dots$, as in [1]. We work with this example since it is mathematically
 212 clear and simple, and is an archetypical example of a chaotic system, being topologically conjugate
 213 to many other well-known systems [17]. However, the analysis could equally apply to more familiar
 214 (continuous) chaotic physical dynamics, such as that of a double pendulum.

215 Let us consider the hypothetical experiment E_k (for fixed $k \geq 1$) which involves iterating the dyadic
 216 map k times (i.e. d^k) on an arbitrary “seed” $\mathbf{x} = x_1x_2\dots$. The outcome of the experiment is then taken
 217 to be the first bit of the resulting sequence $d^k(\mathbf{x}) = x_{k+1}x_{k+2}\dots$, i.e. x_{k+1} . This corresponds to letting
 218 the system evolve for some fixed time k before measuring the result.

219 While the shift d (and hence d^k) is chaotic and generally considered to be unpredictable, it is clearly
 220 simply predictable if we have an extractor that can “see” (or measure) more than k bits of the seed. That
 221 is, take the extractor $\xi_k(\lambda_{\mathbf{x}}) = x_{k+1}$ which clearly extracts only finite information, and the identity Turing
 222 machine I as P_{E_k} so that, for any trial of E_k with parameter $\lambda_{\mathbf{x}}$ we have $P_{E_k}(\xi_k(\lambda_{\mathbf{x}})) = I(x_{k+1}) = x_{k+1}$,
 223 which is precisely the result of the experiment.

224 On the other hand, if we consider that there is some limit l on the “precision” of measurement of \mathbf{x}
 225 that we can perform, the experiment is unpredictable relative to this limited set of extractors Ξ_l defined
 226 such that for every sequence \mathbf{x} and every computable function f there exists λ such that for all $j > l$,
 227 $f(\xi(\lambda)) \neq x_j$. It is clear that for $l = k$, given the limited precision of measurements assumption,
 228 the experiment E_k is unpredictable for Ξ_k . Indeed, if this were not the case, the pair (ξ, P_{E_k}) allowing
 229 prediction would make arbitrarily many correct predictions, thus contradicting the assumption on limited
 230 precision of measurements.

231 This example may appear somewhat artificial, but this is not necessarily so. If one considers the more
 232 physical example of a double pendulum, as mentioned earlier, one can let it evolve for a fixed time t and
 233 attempt to predict its final position (e.g. above or below the horizontal plane) given a set limit l on the
 234 precision of any measurement of the initial position in phase space. If the time t is very short, we may
 235 well succeed, but for long t this becomes unpredictable.

236 This re-emphasises that chaos is an asymptotic property, occurring only strictly at infinite time. While
 237 in the limit it indeed seems to correspond well to unpredictability, in finite time the unpredictability of
 238 chaotic systems is relative: a result of our limits on measurement. Of course, in physical situations such
 239 limits may be rather fundamental: thermal fluctuation or quantum uncertainty seem to pose very real
 240 limits on measurement precision [15], although in most situations the limits actually obtained are of a
 241 far more practical origin.

242 3. Unpredictability in quantum mechanics

243 As we discussed in the introduction, the outcomes of individual quantum measurements are generally
244 regarded as being inherently unpredictable, a fact that plays an important practical role in quantum
245 information theory [18,19]. This unpredictability has many potential origins, of which quantum value
246 indefiniteness is perhaps one of the most promising candidates to be used to certify it more formally.

247 3.1. Quantum value indefiniteness

248 Value indefiniteness is the notion that the outcomes of quantum measurements are not predetermined
249 by any function of the observables and their measurement contexts—that there are no hidden variables.
250 It is thus a formalised notion of indeterminism, and the measurement of such observables results in an
251 outcome not determined before the measurement took place.

252 While it is possible to hypothesise value indefiniteness in quantum mechanics [20], its importance
253 comes from the fact that it can be proven (for systems represented in dimension three or higher Hilbert
254 space) to be true under simple classical hypotheses via the Kochen-Specker theorem [5,21,22]. We will
255 not present the formalism of the Kochen-Specker theorem here, but just emphasise that this gives value
256 indefiniteness a more solid status than a blind hypothesis in the face of a lack of deterministic explanation
257 for quantum phenomena.

258 In [1] we used our model to prove that value indefiniteness can indeed be used to explain quantum
259 unpredictability. Specifically, we showed that *If E is an experiment measuring a quantum value indefinite
260 projection observable, then the outcome of a single trial of E is (simply) unpredictable.*

261 Although value indefiniteness guarantees unpredictability, it relies largely on, and is thus relative to,
262 the Kochen-Specker theorem and its hypotheses [5,21,23], which only holds for systems in three or
263 more dimensional Hilbert space. It is thus useful to know if any other quantum phenomena can be used
264 to certify unpredictability that would be present in two-dimensional systems or in the absence of other
265 Kochen-Specker hypotheses, and if so, what degree of unpredictability is guaranteed.

266 3.2. Complementarity

267 The quantum phenomena of complementarity has also been linked to unpredictability and, contrary to
268 the value indefiniteness pinpointed by the Kochen-Specker theorem, is present in all quantum systems.
269 By itself quantum complementarity is not *a priori* incompatible with value definiteness (there exist
270 automaton and generalised urn models featuring complementarity but not value indefiniteness [24,25])
271 and hence constitutes a weaker hypothesis, even though it is sometimes taken as “evidence” when
272 arguing that value indefiniteness is present in all quantum systems.

273 It is therefore of interest to see if complementarity alone can guarantee some degree of
274 unpredictability, and is an ideal example to apply our model to. This interest is not only theoretical,
275 but also practical as some current quantum random generators [4] operate in two-dimensional Hilbert
276 space where the Kochen-Specker theorem cannot be used to certify value indefiniteness, and would
277 hence seem to (implicitly) rely on complementarity for certification.

278 3.2.1. Quantum complementarity

279 Let us first discuss briefly the notion of quantum complementarity, before we proceed to an analysis
280 of its predictability.

281 The principle of complementarity was originally formulated and promoted by Pauli [26]. It is indeed
282 more of a general principle rather than a formal statement about quantum mechanics, and states that
283 it is impossible to simultaneously measure formally non-commuting observables, and for this reason
284 commutativity is nowadays often synonymous with co-measurability. It is often discussed in the context
285 of the position and momentum observables, but it is equally applicable to any other non-commuting
286 observables such as spin operators corresponding to different directions, such as S_x and S_y , which
287 operate in two-dimensional Hilbert space.

288 Given a pair of such “complementary” observables and a spin- $\frac{1}{2}$ particle, measuring one observable
289 alters the state of the particle so that the measurement of the other observable can no longer be performed
290 on the original state. Such complementarity is closely related to Heisenberg’s original uncertainty
291 principle [27], which postulated that any measurement arrangement for an observable necessarily
292 introduced uncertainty into the value of any complementary observable. For example, an apparatus
293 used to measure the position of a particle, would necessarily introduce uncertainty in the knowledge
294 of the momentum of said particle. This principle and supposed proofs of it have been the subject of
295 longstanding (and ongoing) debate [28–30].

296 More precise are the formal uncertainty relations due to Robertson [31]—confusingly also often
297 referred to as Heisenberg’s uncertainty principle—which state that the standard deviations of the position
298 and momentum observables satisfy $\sigma_x \sigma_p \geq \hbar/2$, and give a more general form for any non-commuting
299 observables A and B . However, this mathematically only places constraints on the variance of repeated
300 measurements of such observables, and does not formally imply that such observables cannot be
301 co-measured, let alone have co-existing definite values, as is regularly claimed [32, Ch. 3].

302 Nonetheless, complementarity is usually taken to mean the stronger statement that it is impossible
303 to simultaneously measure such pairs of observables, and that such measurement of one will result in a
304 loss of information relating to the non-measured observable following the measurement. We will take
305 this as our basis in formalising complementarity, but we do not claim that such a loss of information
306 need be more than epistemic; to deduce more from the uncertainty relations one has to assume quantum
307 indeterminism—that is, value indefiniteness.

308 3.2.2. Complementarity and value definiteness: a toy configuration

309 In order to illustrate that complementarity is not incompatible with value definiteness we briefly
310 consider an example of a toy-model of a system that is value definite but exhibits complementarity. This
311 model was outlined in [25] and concerns a system modelled as an automaton; a different, but equivalent,
312 generalised urn-type model is described in [24].

313 Although this example is just a toy model and does not correspond to a complete quantum system,
314 it represents well many aspects of quantum logic, and serves to show that complementarity itself is not
315 incompatible with value definiteness.

316 The system is modelled as a *Mealy automaton* $\mathcal{A} = (S, I, O, \delta, W)$ where S is the set of states, I and
 317 O the input and output alphabets, respectively, $\delta : S \times I \rightarrow S$ the transition function and $W : S \times I \rightarrow O$
 318 the output function. If one is uncomfortable thinking of a system as an automaton, one can consider the
 319 system as a black-box, whose internal workings as an automaton are hidden. The state of the system
 320 thus corresponds to the state s of the automaton, and each input character $a \in I$ corresponds to a
 321 measurement, the output of which is $W(s, a)$ and the state of the automaton changes to $s' = \delta(s, a)$. To
 322 give a stronger correspondence to the quantum situation, we demand that repeated measurements of the
 323 same character $a \in I$ (i.e. observable) gives the same output: for all $s \in S$ $W(s, a) = W(\delta(s, a), a)$. The
 324 system is clearly value definite, since the output of a measurement is defined prior to any measurement
 325 being made.

326 However, if we have two “measurements” $a, b \in I$ such that $W(s, a) \neq W(\delta(s, b), a)$ then the system
 327 behaves contextually; a and b do not commute. Measuring b changes the state of the system from s to
 328 $s' = \delta(s, b)$, and we lose the ability to know $W(s, a)$.

329 3.3. Complementarity and unpredictability

330 Complementarity tends to be more of a general principle than a formal statement, hence in order
 331 to investigate mathematically the degree of unpredictability that complementarity entails we need to
 332 give complementarity a solid formalism. While several approaches are perhaps possible, following
 333 our previous discussion we choose a fairly strong form of complementarity and consider it not as an
 334 absolute impossibility to simultaneously know the values of non-commuting observables, but rather
 335 as a restriction on our current set of extractors—i.e. using standard quantum measurements and other
 336 techniques we currently have access to.

337 Formally, we say the set of extractors Ξ is *restricted by complementarity* if, for any two incompatible
 338 quantum observables A, B (i.e., $[A, B] \neq 0$), there does not exist an extractor $\xi \in \Xi$ and a computable
 339 function f such that, whenever the value $v(A)$ of the observable A is known², then for all λ , $f(\xi(\lambda)) =$
 340 $v(B)$.

341 This states that, if we know $v(A)$ we have no way of extracting, directly or indirectly, the value $v(B)$
 342 without altering the system. We stress that this doesn't imply that A and B cannot simultaneously have
 343 definite values, simply that we cannot *know* both at once.

² We assume for simplicity that the observables A and B have discrete spectra (as for bounded systems), that is, the eigenvalues are isolated points, and hence the values $v(A)$ and $v(B)$ can be uniquely determined by measurement. Furthermore, since the choice of units is arbitrary (e.g., we can choose $\hbar = 1$) one can generally assume that $v(A)$ and $v(B)$ are rational-valued, and hence can be known ‘exactly’. Even if this were not the case, a finite approximation of $v(A)$ is sufficient to uniquely identify it, and thus is enough here.

For continuous observables it is obviously impossible to identify precisely $v(A)$ or $v(B)$. Such systems are generally idealisations, but one can still handle this case by considering observables A' and B' that measure A and B to some fixed accuracy. Protection by complementarity may depend on this accuracy. For example, for position and momentum, one expects complementarity to apply only when the product of accuracies in position and momentum is less than $\hbar/2$ according to the uncertainty relations.

344 Let us consider an experiment E_C that prepares a system in an arbitrary pure state $|\psi\rangle$, thus giving
 345 $v(P_\psi) = 1$ for the projection observable $P_\psi = |\psi\rangle\langle\psi|$, before performing a projective measurement onto
 346 a state $|\phi\rangle$ with $0 < \langle\psi|\phi\rangle < 1$ (thus $[P_\psi, P_\phi] \neq 0$) and outputting the resulting bit.

347 It is not difficult to see that this experiment is unpredictable relative to an agent whose predicting
 348 power is restricted by complementarity. More formally, if a set of extractors Ξ is restricted by
 349 complementarity, then the experiment E_C described above is unpredictable for Ξ . For otherwise, there
 350 would exist an extractor $\xi \in \Xi$ and a computable predictor P_{E_C} such that, under any repetition procedure
 351 giving parameters $\lambda_1, \lambda_2, \dots$ we have $P_{E_C}(\xi(\lambda_i)) = x_i$ for all i , where x_i is the outcome of the i th
 352 iteration/trial. But if we take $f = P_{E_C}$, then the pair (ξ, f) contradicts the restriction by complementarity,
 353 and hence E_C is unpredictable for Ξ .

354 It is important to note that this result holds regardless of whether the observables measured are value
 355 definite or not, although the value definite case is of more interest. Indeed, if the observables are value
 356 indefinite then we are guaranteed unpredictability without assuming restriction by complementarity, and
 357 hence we gain little extra by considering this situation.

358 As a concrete example, consider the preparation of a spin- $\frac{1}{2}$ particle, for instance an electron, prepared
 359 by in a $S_z = +\hbar/2$ state before measuring the complementary observable $2S_x/\hbar$ producing an outcome
 360 in $\{-1, +1\}$. This could, for example, be implemented by a pair of orthogonally aligned Stern-Gerlach
 361 devices. Next let us assume that the system is indeed value definite. The preparation step means that,
 362 prior to the trial of the experiment being performed, $v(S_z)$ is known, and by assumption $v(S_x)$ exists
 363 (i.e., is value definite) and is thus “contained” in the parameter λ . The assumption that Ξ is restricted
 364 by complementarity means that there is no extractor $\xi \in \Xi$ able to be used by a predictor P_E giving
 365 $P_E(\xi(\lambda_i)) = 2v(S_x)/\hbar = x_i$, thus giving unpredictability for Ξ .

366 As we noted at the start of the section, this is a fairly strong notion of complementarity (although
 367 not the strongest possible). A weaker option would be to consider only that we cannot directly extract
 368 the definite values: that is, there is no $\xi \in \Xi$ such that $\xi(\lambda) = v(S_x)$, for all λ . However, this does not
 369 rule out the possibility that there are other extractors allowing us to indirectly measure the definite values
 370 (unless we take the strong step of assuming Ξ is closed under composition with computable functions, for
 371 example). This weaker notion of complementarity would thus seem insufficient to derive unpredictability
 372 for Ξ , although it would not show predictability either. We would thus, at least for the moment, be left
 373 unsure about the unpredictability of measurements limited by this weak notion of complementarity.

374 4. Unpredictability, computability and complementarity

375 In an effort to try and understand exactly how random quantum randomness—the randomness
 376 generated by measuring unpredictable quantum observables—actually is, we showed in [21] that
 377 quantum value indefiniteness leads to a strong form of incomputability.³ Since this type of
 378 incomputability represents a notion of purely algorithmic unpredictability [1], one may be tempted

³ Technically: A sequence $x_1x_2\dots$ is bi-immune if it contains no computable subsequence, that is, no computable function can compute exactly the values of more than finitely many bits of the sequence.

379 to think that this is a result not so much of quantum value indefiniteness, but rather of quantum
380 unpredictability.

381 In [1], however, we showed that this is not the case: there are unpredictable experiments capable of
382 producing both computable and strongly incomputable sequences when repeated *ad infinitum*. It is thus
383 *a fortiori* true that the same is true for relativised unpredictability, and there is no immediate guarantee
384 that measurements of complementary observables must lead to incomputable sequences as is the case
385 with value indefiniteness.

386 4.1. Incomputability and complementarity

387 Even though the (relativised) unpredictability associated with complementary quantum observables
388 cannot guarantee incomputability, one may ask whether this complementarity may, with reasonable
389 physical assumptions, lead directly to incomputability, much as value indefiniteness does.

390 Here we show this not to be true in the strongest possible way. Specifically, we will show how an,
391 admittedly toy, (value definite) system exhibiting complementarity (and thus unpredictable relative for
392 extractors limited by the complementarity principle) can produce computable sequences when repeated.

393 Consider an experiment E_M involving the prediction of the outcome of measurements on an
394 (unknown) Mealy automaton $M = (Q, \Sigma, \Theta, \delta, \omega)$, which we can idealise as a black box, with $\{x, z\} \in$
395 Σ characters in the input alphabet, output alphabet $\Theta = \{0, 1\}$ and satisfying the condition that x and z
396 are complementary: that is, for all $q \in Q$ we have $\omega(q, z) \neq \omega(\delta(q, x), z)$ and $\omega(q, x) \neq \omega(\delta(q, z), x)$.
397 This automaton is deliberately specified to resemble measurements on a qubit. This very abstract model
398 can be viewed as a toy interpretation of a two-dimensional value definite quantum system, where the
399 outcome of measurements are determined by some unknown, hidden Mealy automaton. Since the
400 Kochen-Specker theorem does not apply to two-dimensional systems, this value definite toy model poses
401 no direct contradiction with quantum mechanics [5], even if it is not intended to be particularly realistic.
402 We complete the specification of E_M by considering a trial of E_M to be the output on the string xz , that
403 is, if the automaton is initially in the state q , the output is $\omega(\delta(q, x), z)$, and the final state is $\delta(\delta(q, x), z)$.
404 This is a clear analogy to the preparation and measurement of a qubit using complementary observables,
405 of the type discussed earlier.

406 Let us show that E_M is unpredictable for a set Ξ_C of extractors that expresses the restriction by
407 complementarity present in Mealy automata. In particular, let us consider the set Ξ_C that, in analogy
408 to the restriction by complementarity of two quantum observables defined earlier, is restricted by an
409 analogue of complementarity for the inputs $x, z \in \Sigma$ in the following sense: *there is no extractor $\xi \in \Xi_C$*
410 *and computable function f such that, if λ_M is the state of a system with Mealy automaton M in a state*
411 *q such that $\delta(q, x) = q$ (or $\delta(q, z) = q$, that is, in an “eigenstate” x or z), then $f(\xi(\lambda_M)) = \omega(q, z)$ (or*
412 *$f(\xi(\lambda_M)) = \omega(q, x)$).* That is, if M is in an “eigenstate” of x , we cannot extract the output of the input
413 z (and similarly for z and x interchanged).

414 Let us assume for the sake of contradiction that E_M is predictable for Ξ_C : that is, there is a predictor
415 P_{E_M} and an extractor $\xi \in \Xi_C$ such that E_M is predictable for Ξ_C . Thus, from the definition of
416 predictability, the pair (P_{E_M}, ξ) must provide infinitely many correct predictions when repeated with the
417 following iteration procedure (in analogy to preparing in an x eigenstate): the black box containing M

418 is prepared by inputting “ x ”, and then the experiment is run and the output recorded. The next repetition
 419 is performed on the same system, preparing the box once again by inputting “ x ” and performing the
 420 experiment. Thus, from the definition of Mealy automata, for each repetition i the automaton M is in
 421 a state q_i such that $\delta(q_i, x) = q_i$ before the i th trial is performed. Thus, the output of the i th trial of
 422 E_M is precisely $\omega(\delta(q_i, x), z) = \omega(q_i, z)$, and for each trial we have $P_{E_M}(\xi(\lambda_i)) = \omega(q_i, z)$, but since
 423 P_{E_M} is a computable function this predictor/extractor pair contradicts the restriction by this form of
 424 complementarity of Ξ_C , and hence we conclude that E_M is unpredictable for Ξ_C .

425 The main question is thus the (in)computability of sequences produced by the concatenation of
 426 outputs from infinite repetitions of E_M . The experiment can be repeated under many different repetition
 427 scenarios, but the simplest is by performing the experiment again on the same black box (and thus with
 428 the same automaton) with the final state of M becoming the initial state for the next repetition⁴. In this
 429 case, the sequence produced is computable—even cyclic—as a result of the automaton M used. Thus,
 430 even if this is not the case under all repetition scenarios, we cannot guarantee that the sequence produced
 431 is incomputable, even though E_M is unpredictable for Ξ_C .

432 We note that one could easily consider slightly more complicated scenarios where the outcomes
 433 are controlled not by a Mealy automaton, but an arbitrary computable—or even, in principle,
 434 incomputable—function; complementarity is agnostic with respect to the computability of the output
 435 of such an experiment. Such a computable sequence may be obviously computable—e.g. $000\dots$, but
 436 it could equally be something far less obvious, such as the digits in the binary expansion of π at prime
 437 indices, e.g. $\pi_2\pi_3\pi_5\pi_7\pi_{11}\dots$. Hence, this scenario cannot be easily ruled out empirically, regardless of
 438 the computability, that is, low complexity, of the resulting sequences. Further emphasising this, we note
 439 that computable sequences can also be Borel-normal, as in Champernowne’s constant or (as conjectured)
 440 π , and thus satisfy many statistical properties one would expect of random sequences.

441 Our point was not to propose this as a realistic physical model—although it perhaps cannot
 442 be dismissed so easily—but to illustrate a conceptual possibility. Value indefiniteness rules this
 443 computability out, but complementarity fails to do the same in spite of its intuitive interpretation as
 444 a form of quantum uncertainty. At best it can be seen as an epistemic uncertainty, as it at least
 445 poses a physical barrier to the knowledge of any definite values. The fact that complementarity
 446 cannot guarantee incomputability is in agreement with the fact that value definite, *contextual* models
 447 of quantum mechanics are perfectly possible [1,33]; such models need not contradict any principle of
 448 complementarity, and can be computable or incomputable.

449 5. Summary

450 In this paper, following on from previous work in [1], we developed a revised and more nuanced
 451 formal model of (un)predictability for physical systems. By considering prediction agents with access
 452 to restricted sets of extractors with which to obtain information for prediction, this model allows various
 453 intermediate degrees of prediction to be formalised.

⁴ Recall that the internal state of M is hidden and not part of the description of E_M , so there is no requirement that it be reset for each repetition.

454 Although models of prediction such as this can be applied to arbitrary physical systems, we have
455 discussed in detail their utility in helping to understand quantum unpredictability, which plays a key role
456 in quantum information and cryptography.

457 We showed that, unlike measurements certified by value indefiniteness, those certified by
458 complementarity alone are not necessarily simply unpredictable: *they are unpredictable relative to the*
459 *ability of the predicting agent to access the values of complementarity observables*—a more epistemic,
460 relativised notion of predictability. This is a general result about complementarity, not specifically
461 in quantum mechanics, and certification by complementarity and value indefiniteness need not be
462 mutually exclusive. Indeed, in dimension three and higher Hilbert space, relative to the assumptions
463 of the Kochen-Specker theorem [21] one has certification by both properties, value indefiniteness thus
464 providing the stronger certification. However, our results are of more importance for two-dimensional
465 systems, since although quantum complementarity is present, this does not necessarily lead to value
466 indefiniteness. While one may postulate value indefiniteness in such cases as well, this constitutes an
467 extra physical assumption, a fact which should not be forgotten [1]. In assessing the randomness of
468 quantum mechanics, one thus needs to take carefully into account all physical assumptions contributing
469 towards the conclusions that one reaches.

470 The fact that quantum complementarity provides a weaker certification than value indefiniteness
471 is emphasised by our final result, showing that complementarity is compatible with the production
472 of computable sequences of bits, something not true for value indefiniteness. Thus, quantum
473 value indefiniteness and the Kochen-Specker theorem appear, for now, essential in certifying the
474 unpredictability and incomputability of quantum randomness.

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477 **Author Contributions**

478 The authors all contributed equally to this paper.

479 **Conflicts of Interest**

480 The authors declare no conflict of interest.

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